# **Exploiting Puzzle Diversity in Puzzle Selection for ESP-like GWAP Systems**

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#### **GWAP = Games with a Purpose**



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PLAYER 1 PLAYER 2



PLAYER 1



PLAYER 2



PLAYER 1



**GUESSING: CAR** 

**GUESSING: HAT** 

**GUESSING: KID** 

PLAYER 2



PLAYER 1



**GUESSING: CAR** 

**GUESSING: HAT** 

**GUESSING: KID** 

PLAYER 2



**GUESSING: BOY** 

**GUESSING: CAR** 

PLAYER 1



**GUESSING: CAR** 

**GUESSING: HAT** 

**GUESSING: KID** 

PLAYER 2



**GUESSING: BOY** 

**GUESSING: CAR** 

Agreement reached: CAR

### Why is it important?

- Some statistics (July 2008)
  - 200,000+ players have contributed 50+ million labels.
  - Each player plays for a total of 91 minutes.
  - The throughput is about 233 labels/player/hour (i.e., one label every 15 seconds)
- Google bought a license to create its own version of the game in 2006 (called Google Image Labeler).





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  - Outsource computational process to human
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- Question: how to evaluate the performance of GWAP systems?





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  - To collect as many labels per puzzle as possible (i.e., quality)
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- Both factors are critical to the performance of the ESP game, but unfortunately they do not complement each other.
- In [14], we formulated the problem as a variant of classic scheduling problems, and proposed an *Optimal Puzzle Selection Algorithm* (OPSA).



# What's the problem?



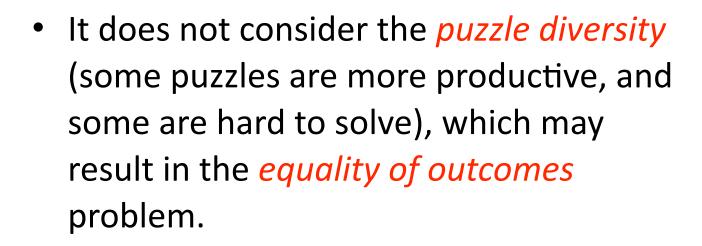
# What's the problem?

 The OPSA scheme determines the optimal number of agreements required for all puzzles based on an analytical model [14].



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#### Contribution

- Using realistic game traces, we identify the puzzle diversity issue in ESP-like GWAP systems.
- We propose the Adaptive Puzzle Selection Algorithm (APSA) to cope with puzzle diversity by promoting equality of opportunity.
- We propose the Weight Sum Tree (WST) to reduce the computational complexity and facilitate the implementation of APSA in real-world systems.
- We show that APSA is more effective than OPSA in terms of the number of agreements reached and the system gain.



### **Adaptive Puzzle Selection Algorithm**

- APSA is inspired by the Additive Increase
   Multiplicative Decrease (AIMD) model of Transmission
   Control Protocol (TCP).
- APSA selects a puzzle to play based on a weighted value  $w_k$ , and the probability that the k-th puzzle will be selected is  $p_k = \frac{w_k}{\sum_{i=1}^K w_i}$

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$$w_k = \begin{cases} 1 & \text{the initial value,} \\ w_k + 1 & \text{if agreements are reached,} \\ \frac{w_k}{2} & \text{if no agreements are reached.} \end{cases}$$



The more productive a puzzle is, the higher probability it will be selected in the next game round.





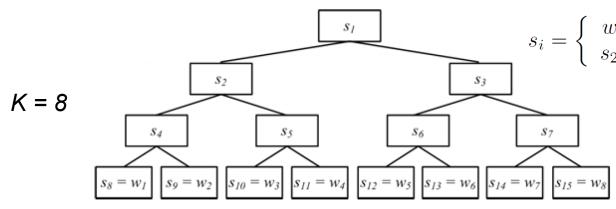
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 $s_i = \begin{cases} w_{i-2^{h-1}+1} & \text{, when } 2^{h-1} \le i < 2^h; \\ s_{2i} + s_{2i+1} & \text{, when } 0 < i < 2^{h-1}. \end{cases}$ 

s<sub>i</sub>: the *i-th* node in the tree h: the height of the tree

- Three cases to maintain the WST
  - After the k-th puzzle is played in a game round
    - Update the w<sub>k</sub> and its ancestors: O(logK)
  - After a puzzle has been removed (say, the k-th puzzle)

 $S_4$ 

 $S_2$ 

 $S_I$ 

 $|s_{10} = w_3| |s_{11} = w_4| |s_{12} = w_5| |s_{13} = w_6|$ 

 $S_3$ 

 $S_6$ 

 $S_7$ 

- Set the  $w_k$  to 0 (to become a virtual puzzle): O(logK)
- After adding a new puzzle (say, the k-th puzzle)
  - Set the w<sub>k</sub> to 1

PReplace the first (leftmost) *virtual* puzzle or rebuild the WST:

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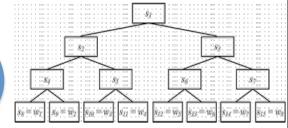
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 Determine a random number r (0 ≤ r ≤ 1), and call the function Puzzle\_Selection(0,r)

**Algorithm 1** The proposed puzzle selection implementation based on the APSA scheme and the weight sum tree data structure.

```
1: Function Puzzle\_Selection(k,r)
```

2: **if** 
$$k \ge 2^{h-1}$$
 **then**

3: Return the 
$$(k-2^{h-1}+1)$$
th puzzle;

4: end if

5: if 
$$r \leq \frac{s_{2k}}{s_1}$$
 then

6: 
$$Puzzle\_Selection(2k, r);$$

7: **else** 

8: 
$$Puzzle\_Selection(2k+1, r-\frac{s_{2k}}{s_1});$$

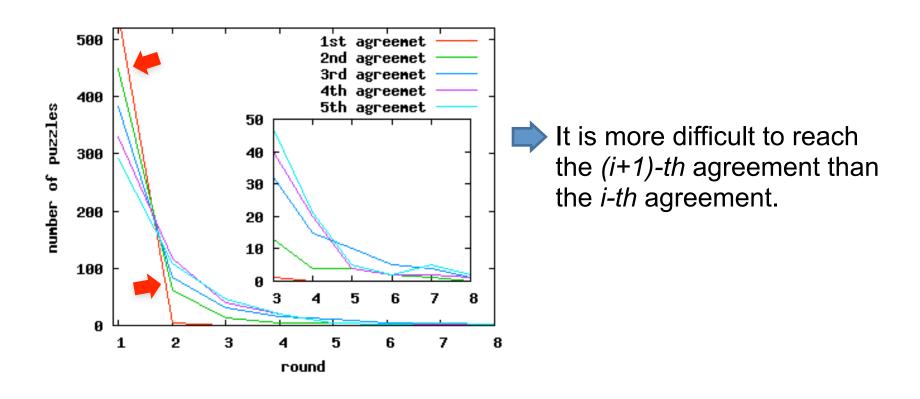
9: end if

#### **Evaluation**

- We evaluated the APSA scheme using trace-based simulations.
- The game trace was collected by the ESP Lite system.
  - The trace was one-month long (from 2009/3/9 to 2009/4/9).
  - The OPSA scheme was used in 1,444 games comprised of 6,326 game rounds. In total, 575 distinct puzzles were played and 3,418 agreements were reached.
  - The dataset is available at:<a href="http://hcomp.iis.sinica.edu.tw/dataset/">http://hcomp.iis.sinica.edu.tw/dataset/</a>

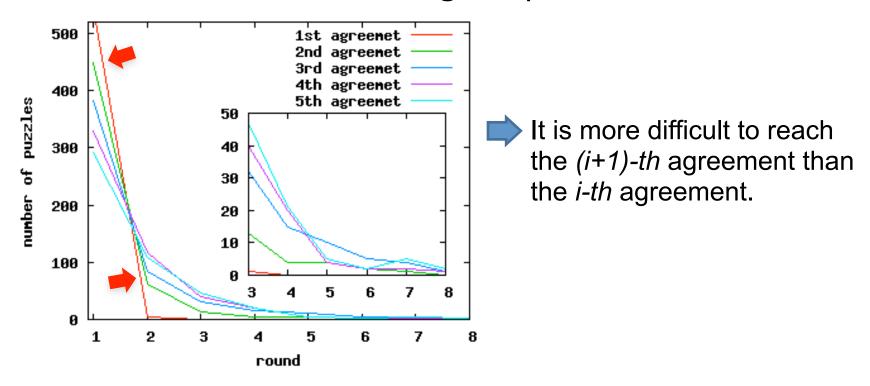


#### **Evaluation – Puzzle Diversity**



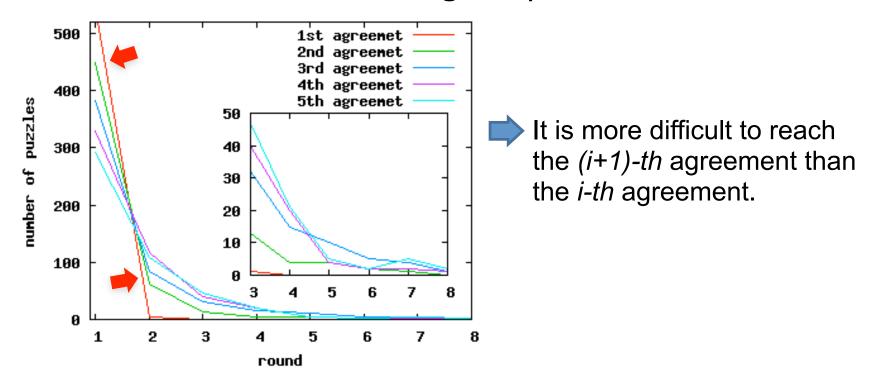
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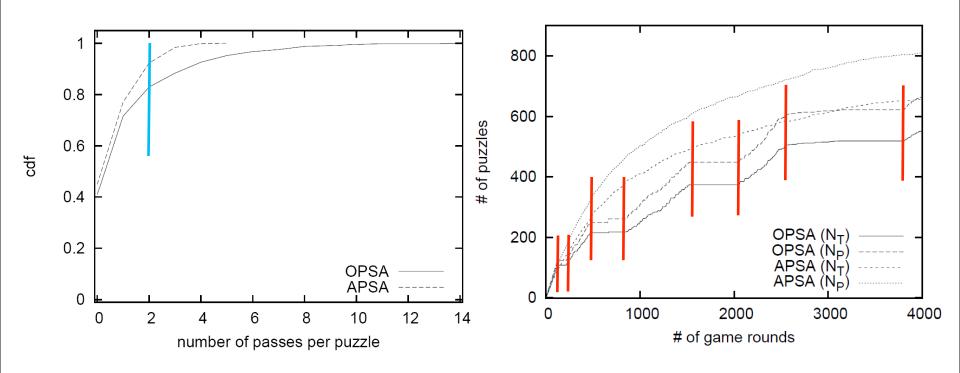
The differences exist among the puzzles.



It is important to consider puzzle diversity!



# **Simulation Results (1)**



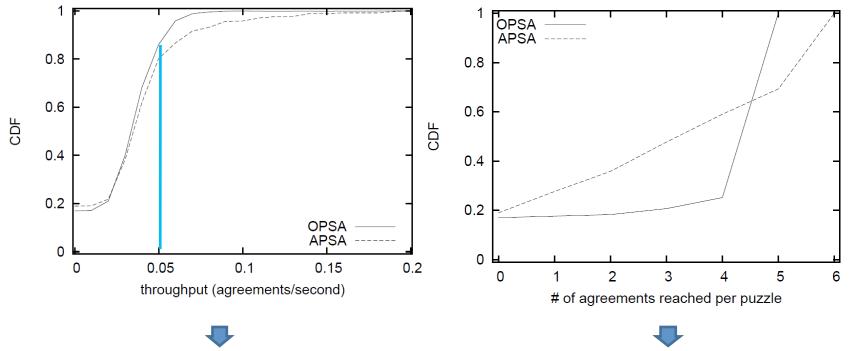


APSA scheme is superior in terms of reducing the number of the passed rounds.

 $N_T$ : # of distinct puzzles with at least one agreement reached  $N_P$ : # of distinct puzzles played



# **Simulation Results (2)**

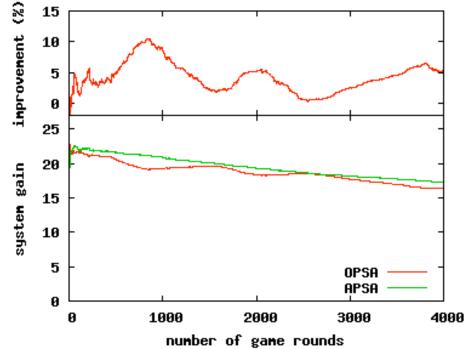


APSA scheme yields more agreements with better per-puzzle throughput

APSA scheme can better accommodate puzzle diversity than the OPSA scheme



#### **System Gain Evaluation**



- APSA always achieves a better system gain than the OPSA scheme (about 5% improvement).
  - The system gain could be improved further by modifying the second part of the metric (e.g., by introducing competition into the system [17]).



#### Summary

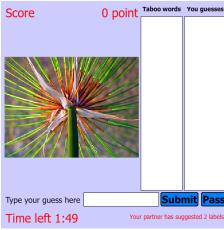
- We identify the puzzle diversity issue in ESP-like GWAP systems.
- We propose the Adaptive Puzzle Selection Algorithm
   (APSA) to consider *individual differences* by promoting equality of opportunity.
- We design a data structure, called Weight Sum Tree
   (WST) to reduce the computational complexity of APSA.
- We evaluate the APSA scheme and show that it is more



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- GWAP API (<a href="http://hcomp.iis.sinica.edu.tw/GWAP\_API/">http://hcomp.iis.sinica.edu.tw/GWAP\_API/</a>)
  - JAVA-based API source codes released
  - ESP Lite: an example of GWAP API
  - ESP Lite dataset (v2010.01.01)

#### **ESP Lite**



http://hcomp.iis.sinica.edu.tw/GWAP/ESPLite

#### **Beauty Ranking**



http://apps.facebook.com/wnranking/

Handsome Ranking



http://apps.facebook.com/menranking/

**Gesture Recognition** 



http://apps.facebook.com/shootit/

#### Babies' Gesture Recognition



http://apps.facebook.com/testforclass/

### **Thank You!**

Network Research Lab: <a href="http://nrl.iis.sinica.edu.tw/">http://nrl.iis.sinica.edu.tw/</a>
GWAP API: <a href="http://hcomp.iis.sinica.edu.tw/GWAP">http://hcomp.iis.sinica.edu.tw/GWAP API/</a>



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