Scout and NegaScout

Tsan-sheng Hsu

徐讚昇

tshsu@iis.sinica.edu.tw

http://www.iis.sinica.edu.tw/~tshsu

Abstract

- It looks like alpha-beta pruning is the best we can do for an **exact generic searching procedure**.
  - What else can be done generically?
  - Alpha-beta pruning follows basically the "intelligent" searching behaviors used by human when domain knowledge is not involved.
  - Can we find some other "intelligent" behaviors used by human during searching?
- Intuition: MAX node.
  - Suppose we know currently we have a way to gain at least 300 points at the first branch.
  - If there is an efficient way to know the second branch is at most gaining 300 points, then there is no need to search the second branch in detail.
    - ▷ *Alpha-beta cut algorithm is one way to make sure of this by returning an exact value.*
    - ▷ *Is there a way to search a tree by only returning a bound?*
    - ▷ *Is searching with a bound faster than searching exactly?*
- Similar intuition holds for a MIN node.

SCOUT procedure

- **It may be possible to verify whether the value of a branch is greater than a value v or not in a way that is faster than knowing its exact value [Judea Pearl 1980].**
- **High level idea:**
  - **While searching a branch T_i of a MAX node, if we have already obtained a lower bound v_ℓ.**
    - ▷ ***First TEST whether it is possible for T_i to return something greater than v_ℓ.***
    - ▷ ***If FALSE, then there is no need to search T_i. This is called fails the test.***
    - ▷ ***If TRUE, then search T_i. This is called passes the test.***
  - **While searching a branch T_j of a MIN node, if we have already obtained an upper bound v_u**
    - ▷ ***First TEST whether it is possible for T_j to return something smaller than v_u.***
    - ▷ ***If FALSE, then there is no need to search T_j. This is called fails the test.***
    - ▷ ***If TRUE, then search T_j. This is called passes the test.***

How to TEST $> v$

procedure TEST$_>$(position p, value v)

// test whether the value of the branch at p is $> v$

- determine the successor positions $p_1, \ldots, p_b$ of p
- if $b = 0$, then // terminal
  - ▷ *if $f(p) > v$ then // f(): evaluating function*
  - ▷ *return TRUE*
  - ▷ *else return FALSE*
- if p is a MAX node, then
  - for $i := 1$ to b do
    - ▷ *if TEST$_>$(p_i, v) is TRUE, then*
      return TRUE // succeed if a branch is $> v$
  - return FALSE // fail only if all branches $\leq v$
- if p is a MIN node, then
  - for $i := 1$ to b do
    - ▷ *if TEST$_>$(p_i, v) is FALSE, then*
      return FALSE // fail if a branch is $\leq v$
  - return TRUE // succeed only if all branches are $> v$

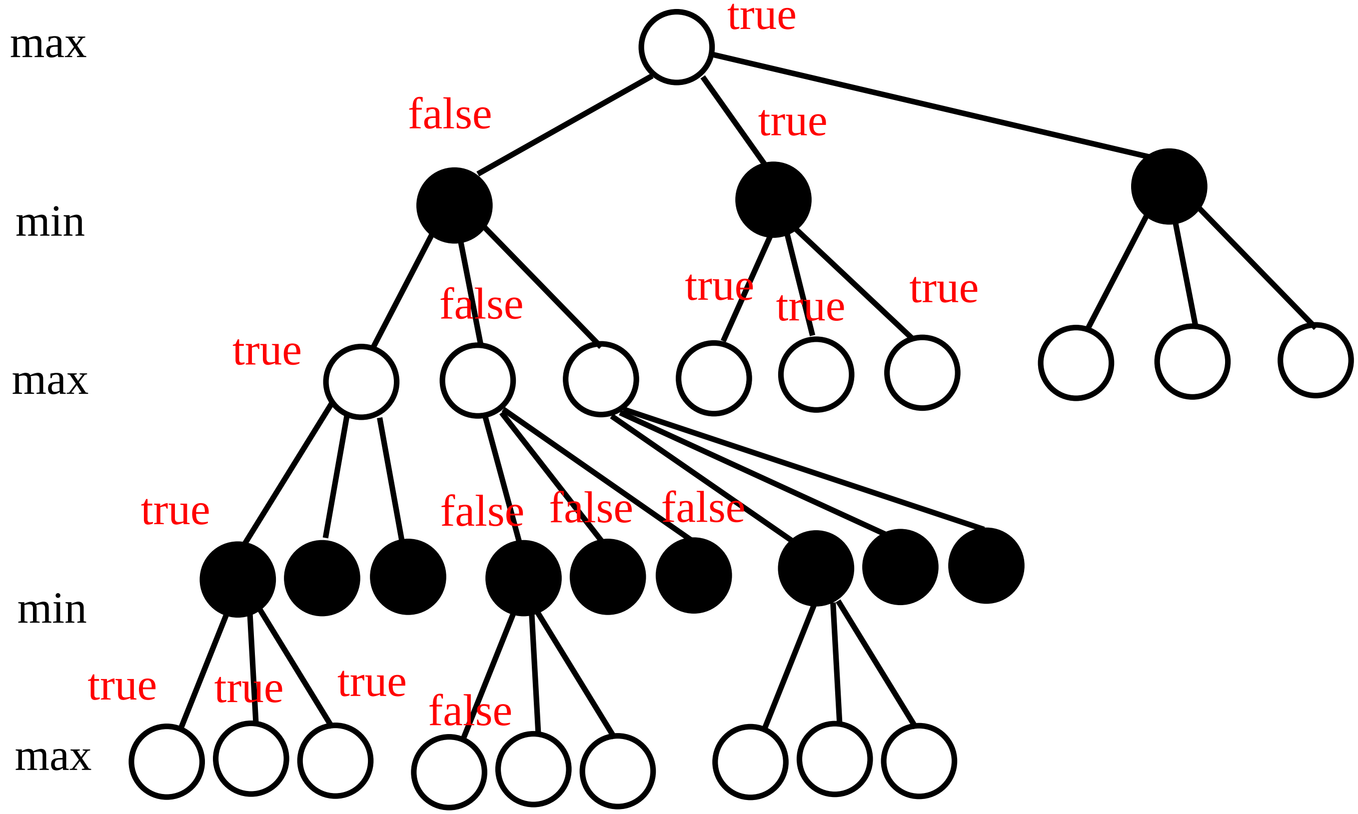
How to TEST $< v$

procedure TEST$_<$(position p, value v)

// test whether the value of the branch at p is $< v$

- **determine the successor positions $p_1, \ldots, p_b$ of p**
- **if $b = 0$, then // terminal**
  - ▷ *if $f(p) < v$ then // f(): evaluating function*
  - ▷ *return TRUE*
  - ▷ *else return FALSE*
- **if p is a MAX node, then**
  - **for $i := 1$ to b do**
    - ▷ *if TEST$_<$(p_i, v) is FALSE, then*
      return FALSE // fail if a branch is $\geq v$
  - **return TRUE // succeed only if all branches $< v$**
- **if p is a MIN node, then**
  - **for $i := 1$ to b do**
    - ▷ *if TEST$_<$(p_i, v) is TRUE, then*
      return TRUE // succeed if a branch is $< v$
  - **return FALSE // fail only if all branches are $\geq v$**

Illustration of TEST$_>$



How to TEST — Discussions

- **Sometimes it may be needed to test for "$\geq v$", or "$\leq v$".**
  - $\boxed{\textbf{TEST}_{>}(p,v) \textbf{ is TRUE}} \equiv \boxed{\textbf{TEST}_{\leq}(p,v) \textbf{ is FALSE}}$
  - $\boxed{\textbf{TEST}_{>}(p,v) \textbf{ is FALSE}} \equiv \boxed{\textbf{TEST}_{\leq}(p,v) \textbf{ is TRUE}}$
  - $\boxed{\textbf{TEST}_{<}(p,v) \textbf{ is TRUE}} \equiv \boxed{\textbf{TEST}_{\geq}(p,v) \textbf{ is FALSE}}$
  - $\boxed{\textbf{TEST}_{<}(p,v) \textbf{ is FALSE}} \equiv \boxed{\textbf{TEST}_{\geq}(p,v) \textbf{ is TRUE}}$
- **Practical consideration:**
  - **Set a depth limit and evaluate the position's value when the limit is reached.**

Main SCOUT procedure

Algorithm SCOUT(position p)

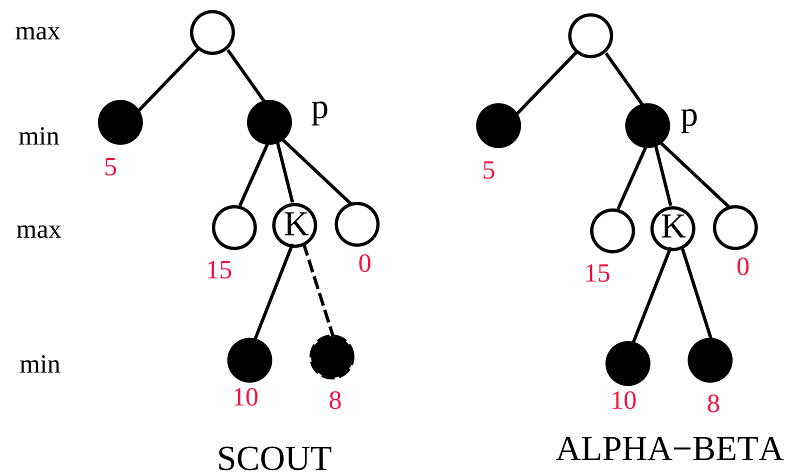
- **determine the successor positions $p_1, \ldots, p_b$**
- **if $b = 0$, then return $f(p)$**
- **else $v = SCOUT(p_1)$ // SCOUT the first branch**
- **if p is a MAX node**
  - **for $i := 2$ to b do**
    - ▷ ***if $TEST_>(p_i, v)$ is TRUE, // TEST first for the rest of the branches***
      then $v = SCOUT(p_i)$ // find the value of this branch if it can be $> v$
- **if p is a MIN node**
  - **for $i := 2$ to b do**
    - ▷ ***if $TEST_<(p_i, v)$ is TRUE, // TEST first for the rest of the branches***
      then $v = SCOUT(p_i)$ // find the value of this branch if it can be $< v$
- **return v**

Discussions for SCOUT (1/3)

- **Note that v is the current best value at any moment.**
- **MAX node:**
  - **For any $i > 1$, if TEST$_>$(p_i, v) is TRUE,**
    - ▷ ***then the value returned by $SCOUT(p_i)$ must be greater than v.***
  - **We say the p_i passes the test if TEST$_>$(p_i, v) is TRUE.**
- **MIN node:**
  - **For any $i > 1$, if TEST$_<$(p_i, v) is TRUE,**
    - ▷ ***then the value returned by $SCOUT(p_i)$ must be smaller than v.***
  - **We say the p_i passes the test if TEST$_<$(p_i, v) is TRUE.**

Discussions for SCOUT (2/3)

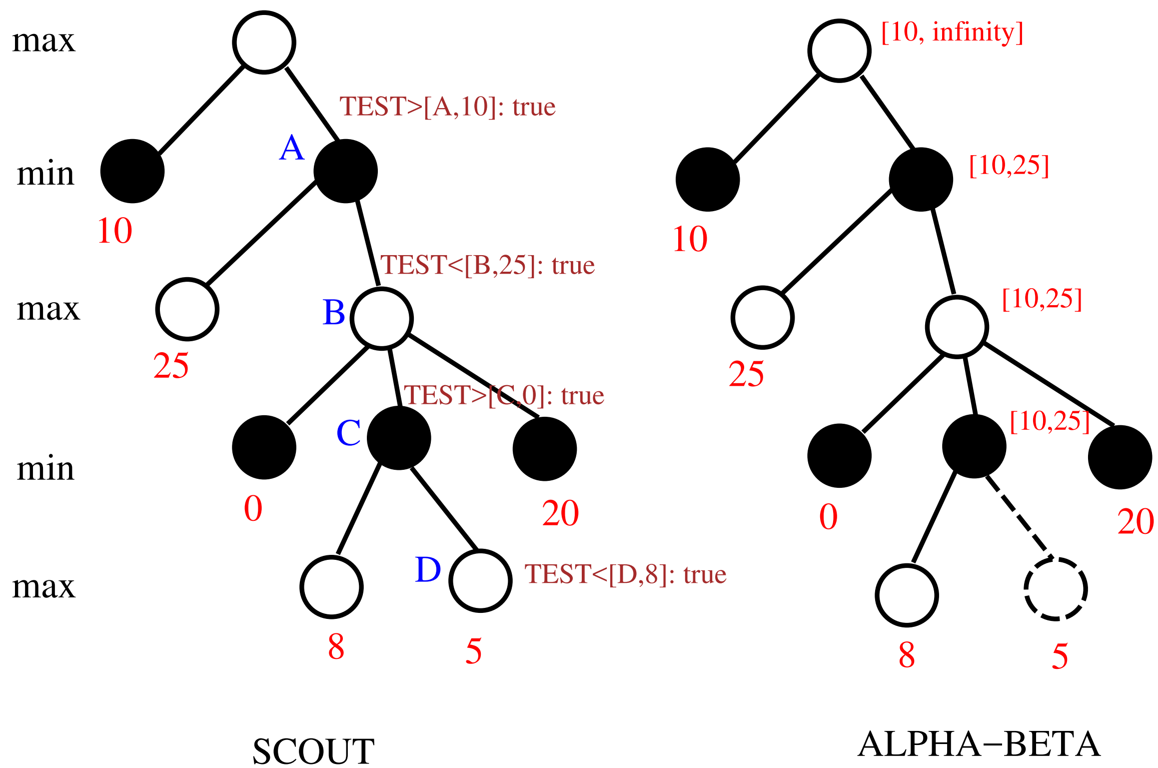
- **TEST which is called by SCOUT may visit less nodes than that of alpha-beta.**



  - **Assume TEST$_>$(p,5) is called by the root after the first branch is evaluated.**
    - ▷ ***It calls TEST$_>$(K,5) which skips K's second branch.***
    - ▷ ***TEST$_>$(p,5) is FALSE, i.e., fails the test, after returning from the 3rd branch.***
    - ▷ ***No need to do SCOUT for the branch p.***
  - **Alpha-beta needs to visit K's second branch.**

Discussions for SCOUT (3/3)

- **SCOUT may pay many visits to a node that is cut off by alpha-beta.**



Number of nodes visited (1/4)

- **For TEST to return TRUE for a subtree T, it needs to evaluate at least**
  - ▷ *one child for a MAX node in T, and*
  - ▷ *and all of the children for a MIN node in T.*
  - ▷ *If T has a fixed branching factor b and uniform depth b, the number of nodes evaluated is $\Omega(b^{\ell/2})$ where ℓ is the depth of the tree.*
- **For TEST to return FALSE for a subtree T, it needs to evaluate at least**
  - ▷ *one child for a MIN node in T, and*
  - ▷ *and all of the children for a MAX node in T.*
  - ▷ *If T has a fixed branching factor b and uniform depth b, the number of nodes evaluated is $\Omega(b^{\ell/2})$.*

Number of nodes visited (2/4)

- **Assumptions:**
  - **Assume a full complete b-ary tree with depth ℓ where ℓ is even.**
  - **The depth of the root, which is a MAX node, is 0.**
- **The total number of nodes in the tree is $\frac{b^{\ell+1}-1}{b-1}$.**
- H_1**: the minimum number of nodes visited by TEST when it returns TRUE.**

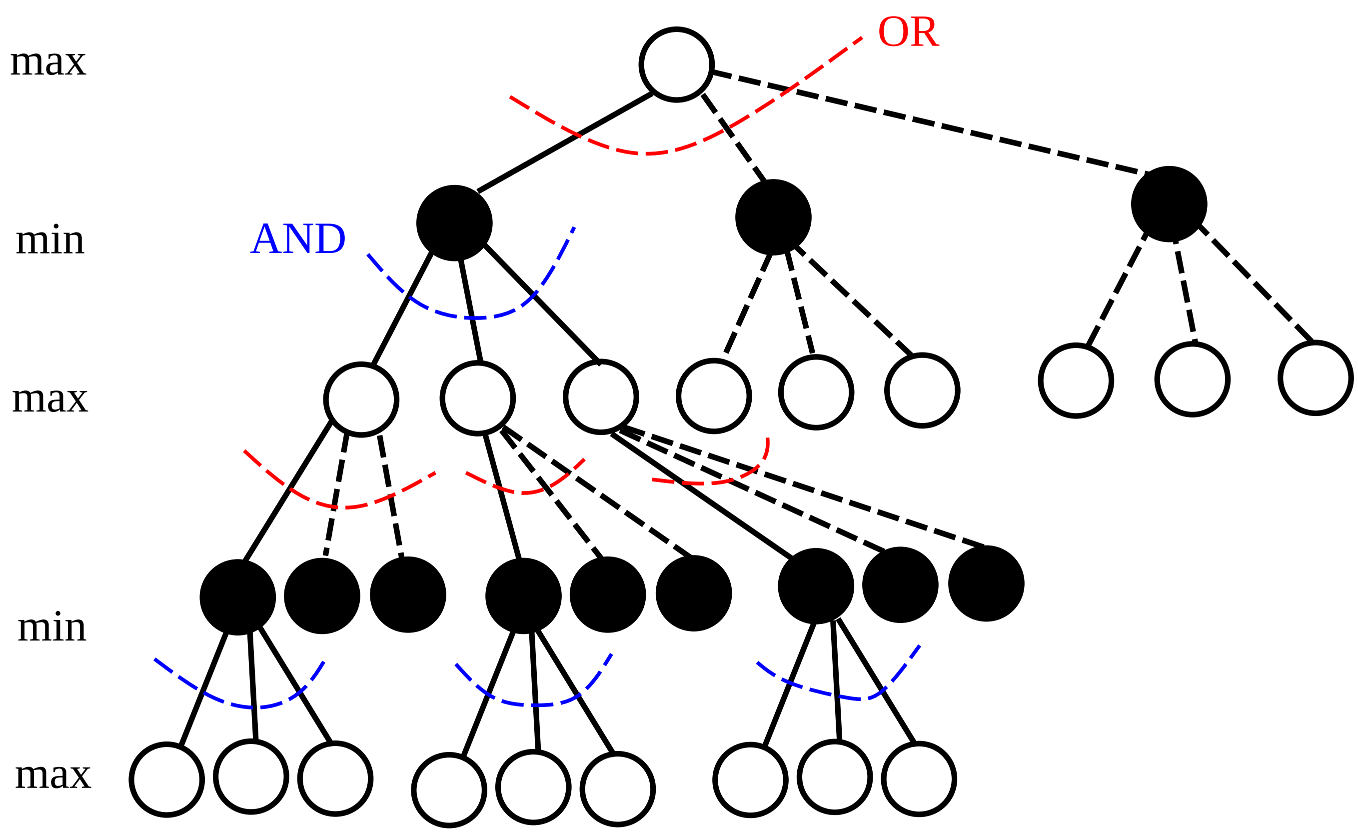
$$
\begin{aligned}
H_1 = & \quad 1 + 1 + b + b + b^2 + b^2 + b^3 + b^3 + \cdots + b^{\ell/2-1} + b^{\ell/2-1} + b^{\ell/2} \\
= & \quad 2 \cdot (b^0 + b^1 + \cdots + b^{\ell/2}) - b^{\ell/2} \\
= & \quad 2 \cdot \frac{b^{\ell/2+1}-1}{b-1} - b^{\ell/2}
\end{aligned}
$$

Number of nodes visited (3/4)

- **Assumptions:**
  - **Assume a full complete b-ary tree with depth ℓ where ℓ is even.**
  - **The depth of the root, which is a MAX node, is 0.**
- H_2**: the minimum number of nodes visited by alpha-beta.**

$$
\begin{aligned}
H_2 &= \sum_{i=0}^{\ell}(b^{\lceil i/2 \rceil} + b^{\lfloor i/2 \rfloor} - 1) \\
&= \sum_{i=0}^{\ell} b^{\lceil i/2 \rceil} + \sum_{i=0}^{\ell} b^{\lfloor i/2 \rfloor} - (\ell + 1) \\
&= \sum_{i=0}^{\ell} b^{\lceil i/2 \rceil} + H_1 - (\ell + 1) \\
&= (1 + b + b + \cdots + b^{\ell/2-1} + b^{\ell/2} + b^{\ell/2}) + H_1 - (\ell + 1) \\
&= (H_1 - 1 + b^{\ell/2} - b^{\ell/2-1}) + H_1 - (\ell + 1) \\
&= 2 \cdot H_1 + b^{\ell/2} - b^{\ell/2-1} - (\ell + 2) \\
&\sim (2.x) \cdot H_1
\end{aligned}
$$

Number of nodes visited (4/4)



Comparisons

- **When the first branch of a node has the best value, then TEST scans the tree fast.**
  - **The best value of the first $i-1$ branches is used to test whether the ith branch needs to be searched exactly.**
  - **If the value of the first $i-1$ branches of the root is better than the value of ith branch, then we do not have to evaluate exactly for the ith branch.**
- **Compared to alpha-beta pruning whose cut off comes from bounds of search windows.**
  - **It is possible to have some cut-off for alpha-beta as long as there are some relative move orderings are "good."**
    - ▷ ***The moving orders of your children and the children of your ancestor who is odd level up decide a cut-off.***
  - **The search bound is updated during the searching.**
    - ▷ ***Sometimes, a deep alpha-beta cut-off occurs because a bound found from your ancestor a distance away.***

Performance of SCOUT (1/2)

- **A node may be visited more than once.**
  - **First visit is to TEST.**
  - **Second visit is to SCOUT.**
    - ▷ *During SCOUT, it may be TESTed with a different value.*
  - **Q: Can information obtained in the first search be used in the second search?**
- **SCOUT is a recursive procedure.**
  - **A node in a branch that is not the first child of a node with a depth of ℓ.**
    - ▷ *Note that the depth of the root is defined to be 0.*
    - ▷ *Every ancestor of you may initiate a TEST to visit you.*
    - ▷ *It can be visited ℓ times by TEST.*

Performance of SCOUT (2/2)

- **Show great improvements on $depth > 3$ for games with small branching factors.**
  - **It traverses most of the nodes without evaluating them preciously.**
  - **Few subtrees remained to be revisited to compute their exact mini-max values.**
- **Experimental data on the game of Kalah show [UCLA Tech Rep UCLA-ENG-80-17, Noe 1980]:**
  - **SCOUT favors "skinny" game trees, that are game trees with high depth-to-width ratios.**
  - **On depth = 5, it saves over 40% of time.**
  - **Maybe bad for games with a large branching factor.**
  - **Move ordering is very important.**
    - ▷ ***The first branch, if is good, offers a great chance of pruning further branches.***

Alpha-beta revisited

- **In an alpha-beta search with a window [$alpha$,$beta$]:**
  - **Failed-high means it returns a value that is larger than or equal to its upper bound $beta$.**
  - **Failed-low means it returns a value that is smaller than or equal to its lower bound $alpha$.**
- **Null or Zero window search:**
  - **Using alpha-beta search with the window [m,$m+1$].**
  - **The result can be either failed-high or failed-low.**
  - **Failed-high means the return value is at least $m+1$.**
    - ▷ ***Equivalent to*** $TEST_>(p,m)$ ***is TRUE.***
  - **Failed-low means the return value is at most m.**
    - ▷ ***Equivalent to*** $TEST_>(p,m)$ ***is FALSE.***
- **The above works for both fail hard and fail soft versions of the alpha-beta algorithm.**

Alpha-Beta + Scout

- **Intuition:**
  - Try to incooperate SCOUT and alpha-beta together.
  - The searching window of alpha-beta if properly set can be used as TEST in SCOUT.
  - Using a searching window is better than using a single bound as in SCOUT.
  - Can also apply alpha-beta cut if it applies.
- **Modifications to the SCOUT algorithm:**
  - Traverse the tree with two bounds as the alpha-beta procedure does.
    - ▷ *A searching window.*
    - ▷ *Use the current best bound to guide the value used in TEST.*
  - Use a fail soft version to get a better result when the returned value is out of the window.

The NegaScout Algorithm – MiniMax (1/2)

- **Algorithm $F4'$(position p, value $alpha$, value $beta$, integer $depth$)**
  - determine the successor positions $p_1, \ldots, p_b$
  - if $b = 0$ // a terminal node
    or $depth = 0$ // $depth$ is the remaining depth to search
    or time is running up // from timing control
    or some other constraints are met // apply heuristic here
  - then return $f(p)$ else
    begin
    - ▷ $m := -\infty$ // m *is the current best lower bound; fail soft*
      $m := \max\{m, G4'(p_1, alpha, beta, depth-1)\}$ // *the first branch*
      if $m \geq beta$ *then return*(m) // *beta cut off*
    - ▷ *for* $i := 2$ *to* b *do*
    - ▷ *9:* $t := G4'(p_i, m, m+1, depth-1)$ // *null window search*
    - ▷ *10:* *if* $t > m$ *then* // *failed-high*
      11: *if* ($depth < 3$ *or* $t \geq beta$)
      12: *then* $m := t$
      13: *else* $m := G4'(p_i, t, beta, depth-1)$ // *re-search*
    - ▷ *14:* *if* $m \geq beta$ *then return*(m) // *beta cut off*

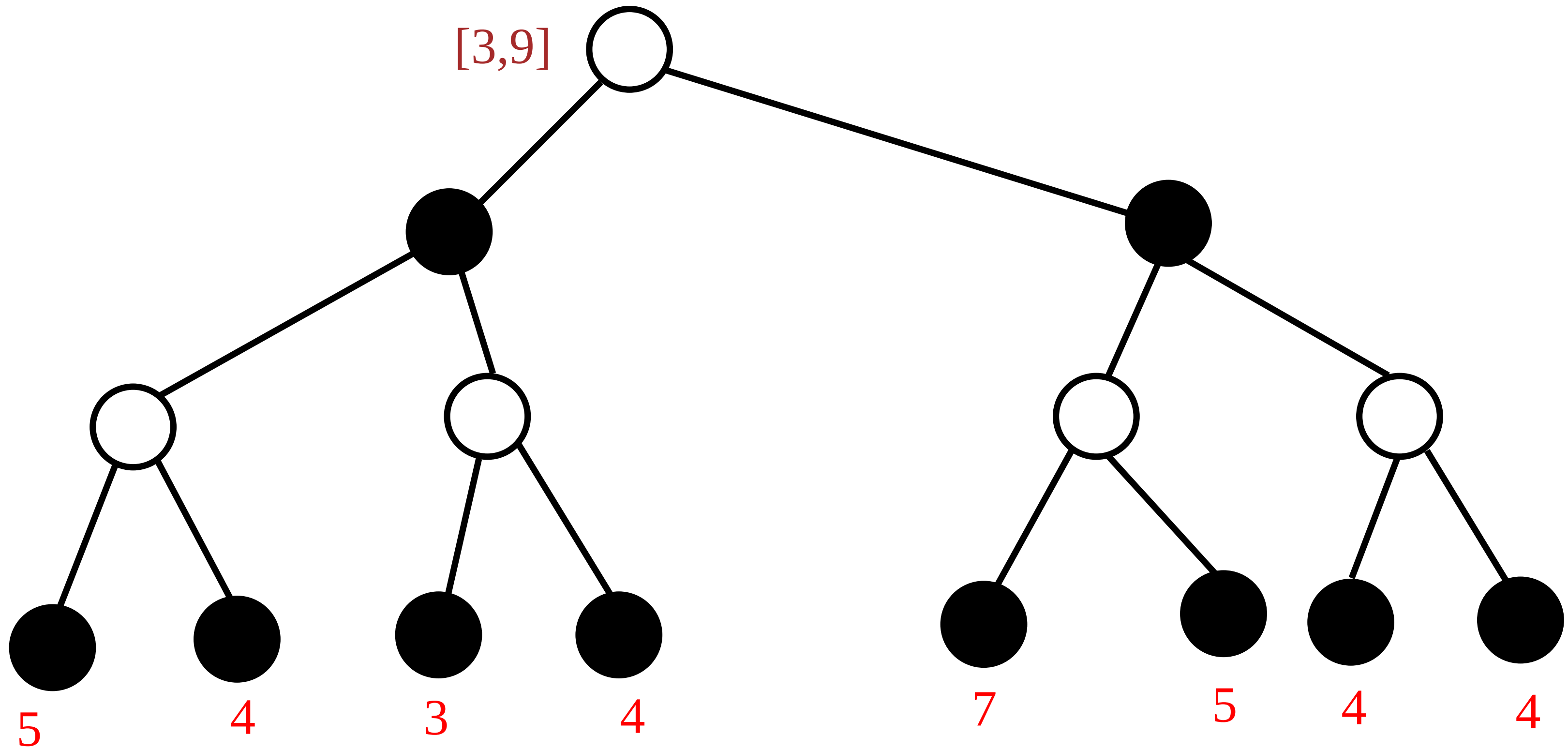
    end
  - return m

The NegaScout Algorithm – MiniMax (2/2)

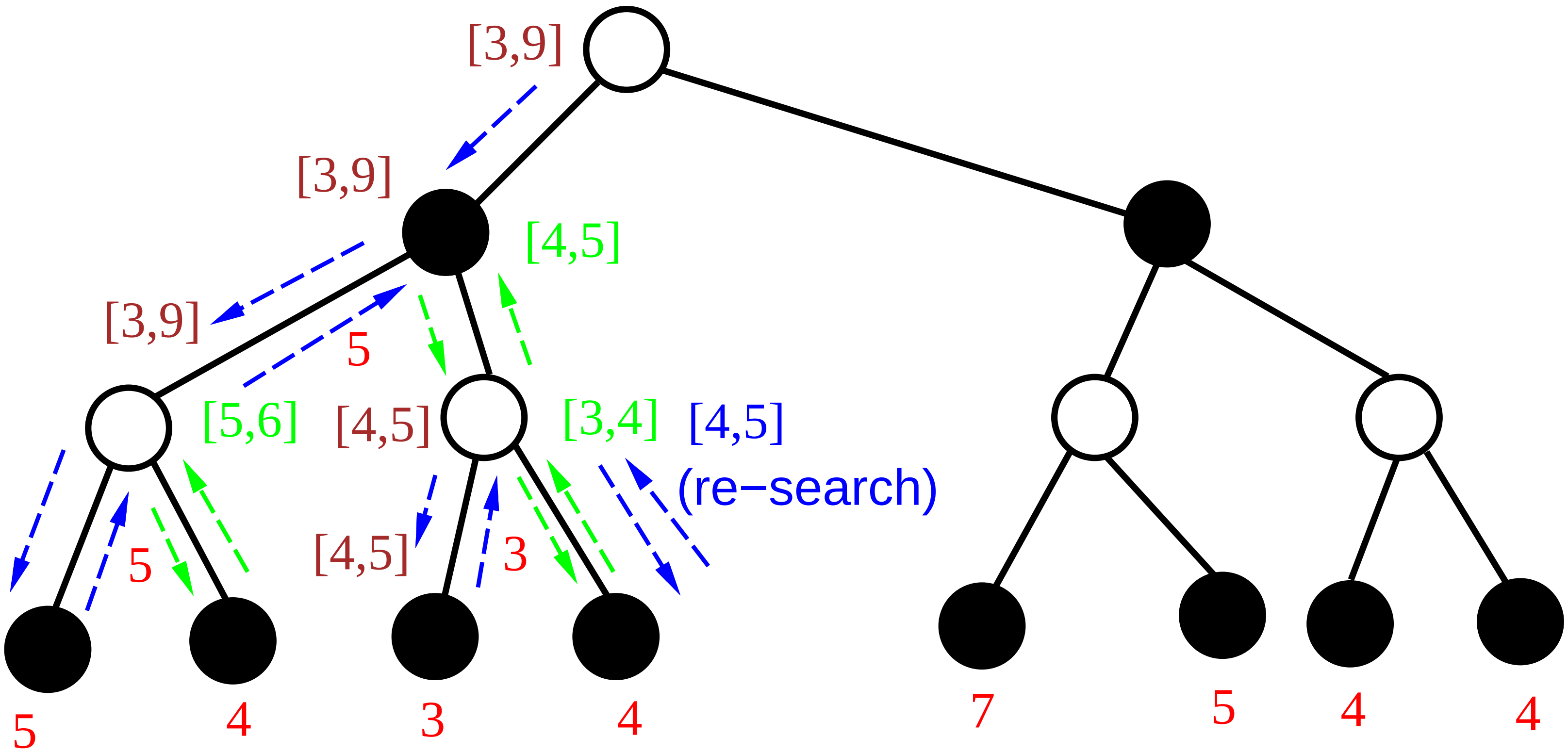
- **Algorithm $G4'$(position p, value $alpha$, value $beta$, integer $depth$)**
  - **determine the successor positions $p_1, \ldots, p_b$**
  - **if $b = 0$ // a terminal node**
    or $depth = 0$ // $depth$ is the remaining depth to search
    or time is running up // from timing control
    or some other constraints are met // apply heuristic here
  - **then return $f(p)$ else**
    begin
    - ▷ $m = \infty$ // m ***is the current best upper bound; fail soft***
      $m := \min\{m, F4'(p_1, alpha, beta, depth - 1)\}$ // ***the first branch***
      if $m \leq alpha$ ***then return***(m) // ***alpha cut off***
    - ▷ ***for*** $i := 2$ ***to*** b ***do***
    - ▷ ***9:*** $t := F4'(p_i, m - 1, m, depth - 1)$ // ***null window search***
    - ▷ ***10:*** ***if*** $t < m$ ***then*** // ***failed-low***
      11: ***if*** $(depth < 3$ ***or*** $t \leq alpha)$
      12: ***then*** $m := t$
      13: ***else*** $m := F4'(p_i, alpha, t, depth - 1)$ // ***re-search***
    - ▷ ***14:*** ***if*** $m \leq alpha$ ***then return***(m) // ***alpha cut off***

    end
  - **return m**

NegaScout – MiniMax version (1/2)



NegaScout – MiniMax version (2/2)



The NegaScout Algorithm

- **Use Nega-MAX format.**
- **Algorithm $F4$(position p, value $alpha$, value $beta$, integer $depth$)**
  - **determine the successor positions $p_1, \ldots, p_b$**
  - **if $b = 0$ // a terminal node**
    or $depth = 0$ //$depth$ is the remaining depth to search
    or time is running up // from timing control
    or some other constraints are met // apply heuristic here
  - **then return $h(p)$ else**
    - ▷ $m := -\infty$ ***// the current lower bound; fail soft***
    - ▷ $n := beta$ ***// the current upper bound***
    - ▷ ***for*** $i := 1$ ***to*** b ***do***
    - ▷ $t := -F4(p_i, -n, -max\{alpha, m\}, depth - 1)$
    - ▷ ***if*** $t > m$ ***then***
      if $(n = beta$ ***or*** $depth < 3$ ***or*** $t \geq beta)$
      then $m := t$
      else $m := -F4(p_i, -beta, -t, depth - 1)$ ***// re-search***
    - ▷ ***if*** $m \geq beta$ ***then return***(m) ***// cut off***
    - ▷ $n := max\{alpha, m\} + 1$ ***// set up a null window***
  - **return m**

Search behaviors (1/3)

- **If the depth is enough or it is a terminal position, then stop searching further.**
  - **Return $h(p)$ as the value computed by an evaluation function.**
  - **Note:**

$$h(p) = \begin{cases} f(p) & \textbf{if depth of } p \textbf{ is 0 or even} \\ -f(p) & \textbf{if depth of } p \textbf{ is odd} \end{cases}$$

- **Fail soft version.**
- **For the first child p_1, search using the normal alpha beta window.**
  - **line 9: normal window for the first child**
  - **the initial value of m is $-\infty$, hence $-max\{alpha, m\} = -alpha$**
    - ▷ *m is the current best value*
  - **that is, searching with the normal window $[alpha, beta]$**

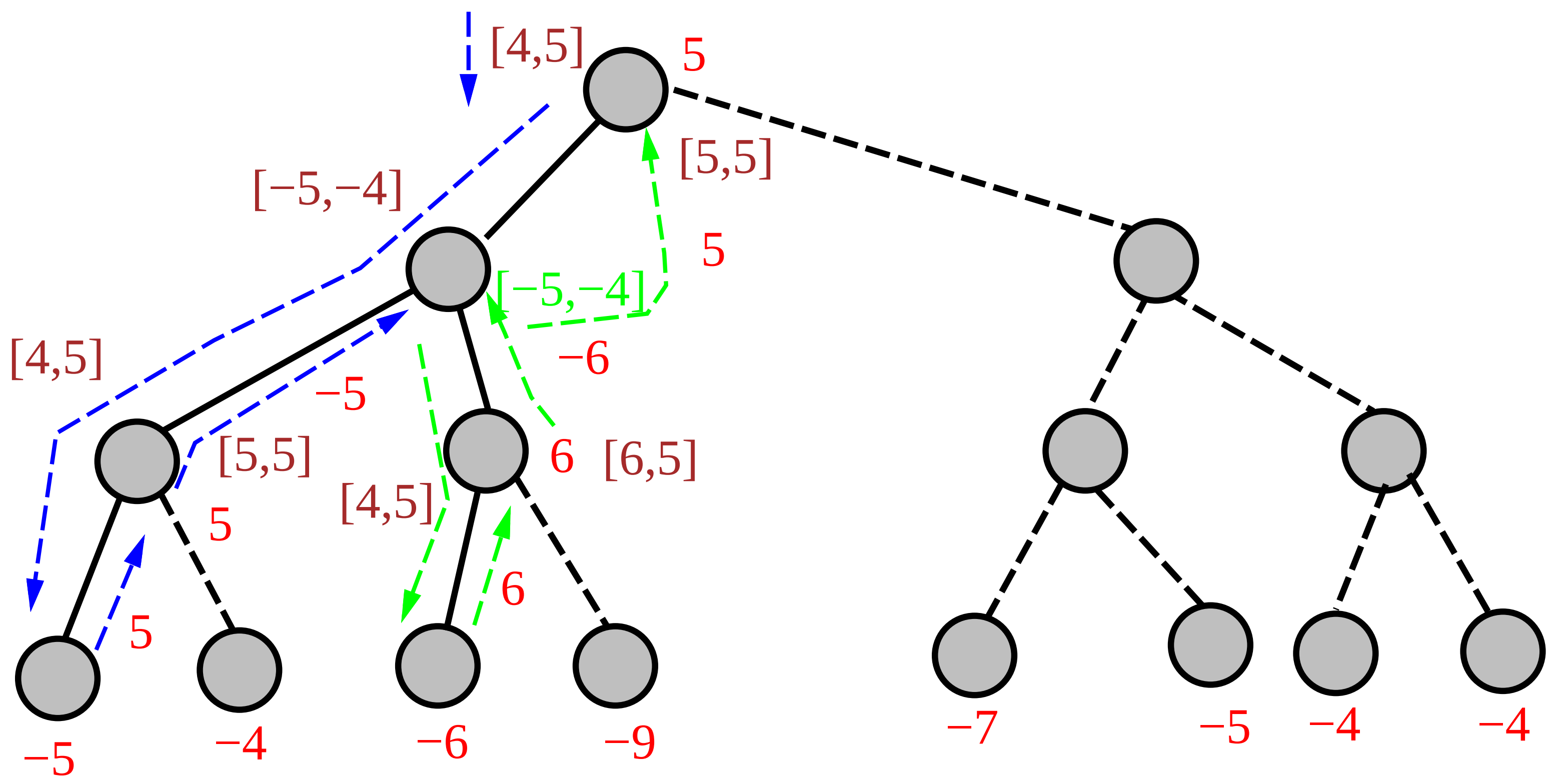
Search behaviors (2/3)

- **For the second child and beyond p_i, $i > 1$, first perform a null window search for testing whether m is the answer.**
  - **line 9: a null-window of $[n-1, n]$ searches for the second child and beyond where $n = max\{alpha, m\} + 1$.**
    - ▷ m ***is best value obtained so far***
    - ▷ $alpha$ ***is the previous lower bound***
    - ▷ m***'s value will be first set at line 12 because*** $n = beta$
    - ▷ ***The value of*** n ***is set at line 15.***
  - **line 11:**
    - ▷ $n = beta$***: we are at the first iteration.***
    - ▷ $depth < 3$***: on a smaller depth subtree, i.e., depth at most 2, NegaScout always returns the best value.***
    - ▷ $t \geq beta$***: we have obtained a good enough value from the failed-soft version to guarantee a beta cut.***

Search behaviors (3/3)

- **For the second child and beyond p_i, $i > 1$, first perform a null window search for testing whether m is the answer.**
  - **line 11: on a smaller depth subtree, i.e., depth at most 2, NegaScout always returns the best value.**
    - ▷ ***Normally, no need to do alpha-beta or any enhancement on very small subtrees.***
    - ▷ ***The overhead is too large on small subtrees.***
  - **line 13: re-search when the null window search fails high.**
    - ▷ ***The value of this subtree is at least t.***
    - ▷ ***This means the best value in this subtree is more than m, the current best value.***
    - ▷ ***This subtree must be re-searched with the the window $[t, beta]$.***
  - **line 14: the normal pruning from alpha-beta.**

Example for NegaScout



Refinements

- **When a subtree is re-searched, it is best to use information on the previous search to speed up the current search.**
  - **Restart from the position that the value t is returned.**
- **Maybe want to re-search using the normal alpha-beta procedure.**
- **$F4$ runs much better with a good move ordering and some form of transposition table.**
  - **Order the moves in a priority list.**
  - **Reduce the number of re-searches.**

Performances

- **Experiments done on a uniform random game tree [Reinefeld 1983].**
  - **Normally superior to alpha-beta when searching game trees with branching factors from 20 to 60.**
  - **Shows about 10 to 20% of improvement.**

Comments

- **Incooperating both SCOUT and alpha-beta.**
- **Used in state-of-the-art game search engines.**
- **The first search, though maybe unsuccessful, can provide useful information in the second search.**
  - **Information can be stored and then reused.**

References and further readings

* **J. Pearl. Asymptotic properties of minimax trees and game-searching procedures.** ***Artificial Intelligence*, 14(2):113–138, 1980.**
* **A. Reinefeld. An improvement of the scout tree search algorithm.** ***ICCA Journal*, 6(4):4–14, 1983.**